# Semidefinite programming in quantum information theory 

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## Linear programming

- Basic task: minimize a linear function of a vector $\vec{x}$ under linear constraints on the elements of $\vec{x}$.
- If there is a solution, it can be solved exactly.


## Semidefinite programming

- Similar, but with semidefinite constraint.
- If there is a solution, it can always be solved exactly.


## Outline

(1) Introduction

- Basic ideas
- Solvable vs. not solvable by SDP
(2) The separability problem
- Separable states
- PPT criterion


## Basic ideas II

- In quantum physics, the density matrix $\varrho$ is a positive semidefinite matrix

$$
\varrho \geq 0
$$

- Its trace is one

$$
\operatorname{Tr}(\varrho)=1
$$

and it is Hermitian

$$
\varrho=\varrho^{\dagger} .
$$

These conditions can easily be included in a semidefinite program.

- When we measure an operator $X$, the expectation value is

$$
\langle X\rangle=\operatorname{Tr}(\varrho X)
$$

## Basic ideas III

- Let us see a simple example. We look for the minimum of

$$
\langle X\rangle=\operatorname{Tr}(\varrho X)
$$

with the condition

$$
\left\langle Y_{n}\right\rangle=\operatorname{Tr}\left(\varrho Y_{n}\right)=y_{n}
$$

for $n=1,2, . ., N$, where $X, Y_{n}$ are operators.

- We optimize over $\varrho$ density matrices.
- This is again doable with semidefinite programming, although, there are better ways to do it.


## Basic ideas III

- Program with MATLAB/YALMIP/MOSEK:

```
X=[001;1 0}]
Y1=[1 0;0 -1];
rho=sdpvar(2,2,'hermitian','complex')
F=[rho>=0]+[trace (rho)==1]+[trace (rho*Y1)==0.2];
    diagnostic=solvesdp(F,trace(X*rho));
    minX=double(trace(X*rho))
```

- Result:

$$
\begin{aligned}
& \min X= \\
&-0.9798
\end{aligned}
$$

## $N$ representability problem I

- Find $\varrho$ of $N$ qudits such that some reduced states are given.
A. J. Coleman, Rev. Mod. Phys. 35, 668 (1963),
for a summary of the literature see in Doherty, Parillo, Spedalieri, PRA 2005.
- Note that if only single-particle reduced states are given, we always have such a $\varrho$.
- If multiparticle reduced states are given, we do not always have such a $\varrho$.


## N representability problem II

- Concrete example: find a two-qubit state such that the reduced states are

$$
\varrho_{1}=\left(\begin{array}{ll}
0.5 & 0.1 \\
0.1 & 0.5
\end{array}\right),
$$

and

$$
\varrho_{2}=\left(\begin{array}{ll}
0.5 & 0.2 \\
0.2 & 0.5
\end{array}\right) .
$$

- The answer: the SDP finds such a state.


## N representability problem III

- Program with MATLAB/YALMIP/MOSEK:
\% 2 qubits
rho=sdpvar (4,4,'hermitian','complex')
\% Reduced states
rhol=[0.5 0.1;0.1 0.5];
rho2=[0.5 0.2; 0.2 0.5];
$\mathrm{F}=[$ rho $>=0]+[$ trace $($ rho $)==1]$;
\% reduced states using an external routine $\mathrm{F}=\mathrm{F}+[\mathrm{keep}$ _nonorm (rho, 1) = = rhol];
$\mathrm{F}=\mathrm{F}+[\mathrm{keep}$ _nonorm (rho, 2) ==rho2];
diagnostic=solvesdp ( $\mathrm{F}, 0$ ) ;
is_there_a_problem=diagnostic.problem
rho_solution=double (rho)


## N representability problem III

- Result

$$
\begin{aligned}
& \text { is_there_a_problem = } \\
& \begin{array}{rlll}
0 & & \\
\text { rho_solution }= & & & \\
0.2500 & 0.0500 & 0.1000 & 0.0344 \\
0.0500 & 0.2500 & 0.0344 & 0.1000 \\
0.1000 & 0.0344 & 0.2500 & 0.0500 \\
0.0344 & 0.1000 & 0.0500 & 0.2500
\end{array}
\end{aligned}
$$

## $N$ representability problem IV

- Concrete example: find a two-qubit state such that the reduced states are

$$
\varrho_{1}=\left(\begin{array}{ll}
0.5 & 0.5 \\
0.5 & 0.5
\end{array}\right),
$$

and

$$
\varrho_{2}=\left(\begin{array}{ll}
0.5 & 0.5 \\
0.5 & 0.5
\end{array}\right) .
$$

- Note that the states are pure states and they are the eigestates of $\sigma_{x}$ with an eigenvalue +1 .
- The answer: the SDP finds such a state, it will find $\varrho_{1} \otimes \varrho_{2}$.


## $N$ representability problem V

- Program with MATLAB/YALMIP/MOSEK:
\% 2 qubits
rho=sdpvar (4,4,'hermitian','complex')
\% Reduced states
rhol=[0.5 0.5;0.5 0.5];
rho2 $=[0.5$ 0.5; 0.5 0.5];
$\mathrm{F}=[$ rho $>=0]+[$ trace $($ rho $)==1]$;
\% reduced states using an external routine $\mathrm{F}=\mathrm{F}+\left[\mathrm{keep} \_\right.$nonorm (rho, 1) ==rho1];
$\mathrm{F}=\mathrm{F}+$ [keep_nonorm (rho, 2) ==rho2];
diagnostic=solvesdp ( $\mathrm{F}, 0$ ) ;
is_there_a_problem=diagnostic.problem rho_solution=double (rho)


## $N$ representability problem VI

- Result

$$
\begin{array}{lll}
\text { is_there_a_problem }= \\
0 & \\
\text { rho_solution }= & & \\
0.2500 & 0.2500 & 0.2500
\end{array} 00.2500
$$

## Nonlinear optimization

- Let us see a simple example. We look for the minimum of

$$
\left\langle X_{1}\right\rangle^{2}+\left\langle X_{2}\right\rangle^{2}=\operatorname{Tr}\left(\varrho X_{1}\right)^{2}+\operatorname{Tr}\left(\varrho X_{2}\right)^{2}
$$

with the condition

$$
\left\langle Y_{n}\right\rangle=\operatorname{Tr}\left(\varrho Y_{n}\right)=y_{n}
$$

for $n=1,2, . ., N$.

- We optimize over $\varrho$ density matrices.
- This is again doable with semidefinite programming, minimizing $t_{1}+t_{2}$ using the constraints

$$
\left(\begin{array}{cc}
t_{k} & \operatorname{Tr}\left(\varrho X_{k}\right) \\
\operatorname{Tr}\left(\varrho X_{k}\right) & 1
\end{array}\right) \geq 0
$$

for $k=1,2$.

## Nonlinear optimization II

- Concrete example: we minimize

$$
\left\langle\sigma_{x}\right\rangle^{2}+\left\langle\sigma_{z}\right\rangle^{2}
$$

with the constraint

$$
\left\langle Y_{1}\right\rangle=0.2
$$

where

$$
Y_{1}=\left(\begin{array}{cc}
1 & i \\
-i & 1
\end{array}\right)
$$

- The answer: we find a density matrix that corresponds to the minimum. For that state, we have

$$
\left\langle\sigma_{x}\right\rangle^{2}+\left\langle\sigma_{z}\right\rangle^{2}=0.01
$$

## Nonlinear optimization III

- Program with MATLAB/YALMIP/MOSEK:

```
X1=[0 1;1 0]; X2=[11 0;0 -1]; Y1=[[1 i;-i -1];
rho=sdpvar(2,2,'hermitian','complex')
t=sdpvar(1,1,'full','real')
F=[rho>=0]+[trace(rho)==1];
F=F+[trace(Y1*rho)==0.2];
M=[t trace(X1*rho);trace(X1*rho) t];
F}=\textrm{F}+[\textrm{M}>=0]
    diagnostic=solvesdp(F,t);
    is_there_a_problem=diagnostic.problem
    rho_solution=double(rho)
```


## Nonlinear optimization IV

- Results:

```
is_there_a_problem =
    0
rho_solution =
    0.5500 + 0.0000i 0.0000 + 0.0500i
    0.0000 - 0.0500i 0.4500 + 0.0000i
>> trace(X2*rho_solution)
ans =
    0.1000
>> trace(X1*rho_solution)
ans =
    0
```


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## Solvable vs. not solvable by SDP

- Thus, we can minimize a convex function over the convex set of density matrices.

- However, we cannot maximize a function over the convex set of density matrices efficiently - the maximum is taken at the boundaries.



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## Mixed states: separable states vs. entangled states

- For the mixed case, the definition of a separable state is (Werner 1989)

$$
\rho_{\mathrm{sep}}=\sum_{k} p_{k}\left[\rho_{k}^{(1)}\right]_{A} \otimes\left[\rho_{k}^{(2)}\right]_{B} .
$$

A state that is not separable, is entangled.

- It is not possible to create entangled states from separable states, with LOCC.
- From many copies of two-qubit mixed entangled states, we can always distill a singlet using Local Operations and Classical Communication (LOCC).
- This is not true for higher dimensional systems. Not all quantum states are distillable.


## Convex sets

## Entangled states

Separable states

## Bipartite systems I

- Naive question: can we decide whether a state is separable with SDP? No, because we would need a constraint of the type

$$
\varrho=\left(\varrho_{1}\right)_{A} \otimes\left(\varrho_{2}\right)_{B}
$$

- Alternatively, we would need a constraint for the reduced states of the nth subsytem

$$
\operatorname{Tr}\left(\varrho_{\mathrm{red}, n}^{2}\right)=1
$$

## Bipartite systems II

- How can we check separability using a brute force method? We can look for a separable decomposition with some $\rho_{k}^{(1)}, \rho_{k}^{(2)}$ numerically.
- Simpler problem, maximum for an operator expectation value for separable states

$$
\left.\max _{\rho_{\text {sep }}} \operatorname{Tr}\left(X \rho_{\text {sep }}\right)=\max _{\Psi_{1}, \Psi_{2}}\left\langle\Psi_{1} \mid\left\langle\Psi_{2}\right| X \mid \Psi_{2}\right\rangle \Psi_{1}\right\rangle
$$

- Numerically, we can try to find the maximum. In practice, we will find the maximum or something lower.


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## The positivity of the partial transpose (PPT) criterion

## Definition

For a separable state $\varrho$ living in $A B$, the partial transpose is always positive semidefinite

$$
\varrho^{T A} \geq 0
$$

If a state does not have a positive semidefinite partial transpose, then it is entangled. A. Peres, PRL 1996; Horodecki etal., PLA 1997.

- Partial transpose means transposing according to one of the two subsystems.
- For separable states

$$
(T \otimes \mathbb{1}) \varrho=\varrho^{T A}=\sum_{k} p_{k}\left(\varrho_{k}^{(1)}\right)^{T} \otimes \varrho_{k}^{(2)} \geq 0
$$

## The positivity of the partial transpose (PPT) criterion II

- How to obtain the partial transpose of a general density matrix? Example: $3 \times 3$ case.



# The positivity of the partial transpose (PPT) criterion III 

- If the relation

$$
\varrho^{T A} \geq 0
$$

is violated then the state is entangled!

- For $2 \times 2$ and $2 \times 3$ systems it detects all entangled states.
- For larger systems, there are entangled states for which

$$
\varrho^{T A} \geq 0
$$

hold. They are bound entangled, not distillable.

## Convex sets

## non-PPT Entangled states

## PPT Entangled states

Separable states

## The positivity of the partial transpose (PPT) criterion IV

- Semidefinite programming can be used to optimize over PPT states.
- Find the minimum of an operator expectation value for PPT states:

Minimize

$$
\langle X\rangle_{\varrho} \equiv \operatorname{Tr}(X \varrho)
$$

such that

$$
\begin{aligned}
\varrho & =\varrho^{\dagger} \\
\varrho & \geq 0 \\
\varrho^{T A} & \geq 0 \\
\operatorname{Tr}(\varrho) & =1
\end{aligned}
$$

# The positivity of the partial transpose (PPT) criterion V 

- Concrete example: look for the minimum of

$$
\left\langle\sigma_{x} \otimes \sigma_{x}+\sigma_{y} \otimes \sigma_{y}\right\rangle
$$

for PPT states.

- The answer: the minimum is -1 .


## The positivity of the partial transpose (PPT) criterion VI

- Program with MATLAB/YALMIP/MOSEK:

```
sigmax=[0 1;1 0];sigmay=[0 -i;i 0];
% We want to minimze <A>
A=kron(sigmax,sigmax) +kron(sigmay,sigmay);
% Two qubits
rho=sdpvar(4,4,'hermitian',' complex')
F=[rho>=0]+[trace(rho)==1];
% Using an external partial transpose routine
F=F+[pt_nonorm(rho,1)>=0];
diagnostic=solvesdp(F,trace(A*rho));
minimum=double(trace(A*rho))
```


## The positivity of the partial transpose (PPT) criterion VII

- Result:

$$
\begin{aligned}
& \text { minimum }= \\
& -1.0000
\end{aligned}
$$

## The positivity of the partial transpose (PPT) criterion VIII

- Thus, we find that the mininum of

$$
\left\langle\sigma_{x} \otimes \sigma_{x}+\sigma_{y} \otimes \sigma_{y}\right\rangle
$$

for PPT states is -1 .

- This is like finding a lower bound on the minimum for separable states.
- In practice, we often find the minimum for separable states, as in the example above.
G. Tóth, W. Wieczorek, R. Krischek, N. Kiesel, P. Michelberger, and H. Weinfurter, New J. Phys. 2009.


## The positivity of the partial transpose (PPT) criterion IX

- We can ask: is there a PPT fulfilling certain constraints?

Look for $\varrho$ such that

$$
\begin{aligned}
\varrho & =\varrho^{\dagger}, \\
\varrho & \geq 0 \\
\varrho^{T A} & \geq 0, \\
\operatorname{Tr}(\varrho) & =1, \\
\operatorname{Tr}\left(X_{k} \varrho\right) & =x_{k} \text { for } k=1,2, \ldots, K .
\end{aligned}
$$

- If there is not such a $\varrho$ then the state fulfilling the constraints is not PPT, and it is entangled (or it is not physical).
- One can use this to detect entanglement in experiments.


# The positivity of the partial transpose (PPT) criterion X 

- Concrete example: is there a PPT state with

$$
\left\langle\sigma_{x} \otimes \sigma_{x}+\sigma_{y} \otimes \sigma_{y}\right\rangle=1.1 ?
$$

- The answer: there is no such a PPT state.


## The positivity of the partial transpose (PPT) criterion XI

- Program with MATLAB/YALMIP/MOSEK: CHANGE!!!

```
sigmax=[0 1;1 0];sigmay=[0 -i;i 0];
A=kron(sigmax,sigmax)+kron(sigmay,sigmay);
```

\% Two qubits
rho=sdpvar(4,4,'hermitian',' complex')
$\mathrm{F}=[$ rho>=0] + [trace (rho) ==1];
\% Using an external partial transpose routine
\% Condition $<A>=1.1$
$\mathrm{F}=\mathrm{F}+[\mathrm{pt}$ _nonorm (rho, 1) $>=0]+[$ trace $(\mathrm{A} *$ rho) $==1.1]$;
diagnostic=solvesdp(F,0);
is_there_a_problem=diagnostic.problem

# The positivity of the partial transpose (PPT) criterion XII 

- Result:

$$
\begin{gathered}
\text { is_there_a_problem }= \\
1
\end{gathered}
$$

- That is, there is not such a PPT quantum state.

